Unsolvable Problems

Part One

Outline for Today

- ***Self-Reference Revisited***
  - Programs that compute on themselves.
- ***Self-Defeating Objects***
  - Objects “too powerful” to exist.
- ***The Fortune Teller***
  - Can you escape the future?
- ***Why Do Programs Loop?***
  - … and can we eliminate loops?
- ***Undecidable Problems***
  - Something beyond the reach of algorithms.

Recap from Last Time

R and **RE**

- A language L is ***recognizable*** if there is a TM M with the following property:

$$\forall w \in \Sigma^*. \ (M \text{ accepts } w \leftrightarrow w \in L).$$

- That is, for any string w:
  - If $w \in L$, then M accepts w.
  - If $w \notin L$, them M does not accept w.
    - It might reject w, or it might loop on w.
- This is a "weak" notion of solving a problem.
- The class **RE** consists of all the recognizable languages.

R and **RE**

- A language L is ***decidable*** if there is a TM M with the following properties:

$$\forall w \in \Sigma^*. (M \text{ accepts } w \leftrightarrow w \in L).$$

***M* halts on all inputs.**

- That is, for any string w:
  - If $w \in L$, then M accepts w.
  - If $w \notin L$, then M rejects w.
- This is a "strong" notion of solving a problem.
- The class **R** consists of all the decidable languages.

The Universal TM

- The ***universal Turing machine***, denoted U_{TM}, is a TM with the following behavior: when run on a string $\langle M, w\rangle$, where M is a TM and w is a string, U_{TM} will

  … accept $\langle M, w\rangle$ if M accepts w,

  … reject $\langle M, w\rangle$ if M rejects w, and

  … loop on $\langle M, w\rangle$ if M loops on w.

- $\mathbf{A_{TM}}$ is the language recognized by the universal TM. This is the language

$$\mathbf{A_{TM} = \{ \langle \mathit{M}, \mathit{w} \rangle \mid \mathit{M} \text{ is a TM and } \mathit{M} \text{ accepts } \mathit{w} \}}$$

New Stuff!

Part One: Self-Defeating Objects

A ***self-defeating object*** is an object whose essential properties ensure it doesn't exist.

Question: Why is there no largest integer?

Answer: Because if n is the largest integer, what happens when we look at $n+1$?

Self-Defeating Objects

Theorem: There is no largest integer.

Proof sketch: Suppose for the sake of contradiction that there is a largest integer. Call that integer n.

Consider the integer $n+1$.

Notice that $n < n+1$.

But then n isn't the largest integer.

Contradiction! ■-*ish*

Self-Defeating Objects

Theorem: There is no largest integer.

Proof sketch: Suppose for the sake of contradiction that there is a largest integer. Call that integer n.

Consider the integer $n+1$.

Notice that $n < n+1$.

But then n isn't the larg

Contradiction! ■-*ish*

We're using n to construct something that undermines n, hence the term "self-defeating."

An Important Detail

Careful – we're assuming what we're trying to prove!

Claim: There is a largest integer.

Proof: Assume x is the largest integer.

Notice that $x > x - 1$.

So there's no contradiction. ■-*ish*

How do we know there's no contradiction? We just checked one case.

Self-Defeating Objects

- If you can show

$$\boldsymbol{x \text{ exists} \to \bot}$$

  then you know that *x* doesn't exist. (This is a proof by contradiction.)

- If you can show

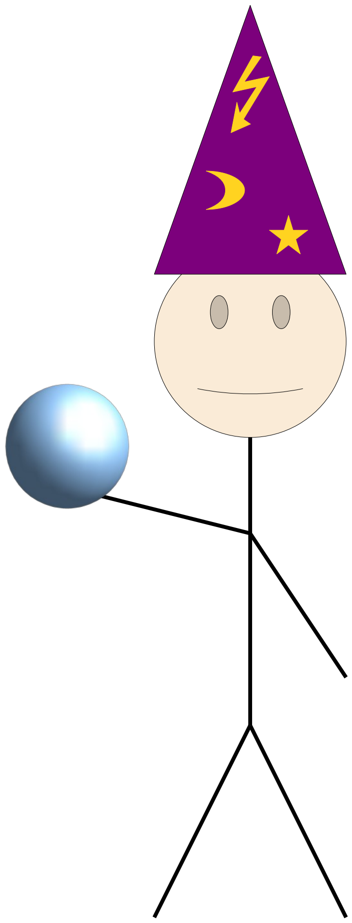
$$\boldsymbol{x \text{ exists} \to \top}$$

  you cannot conclude that *x* exists. (This is not a valid proof technique.)

Part Two: The Fortune Teller

The Fortune Teller

- A fortune teller appears who claims they can see into anyone’s future.
- For a nominal fee, the fortune teller will tell you anything you want to know about the future.



The Fortune Teller

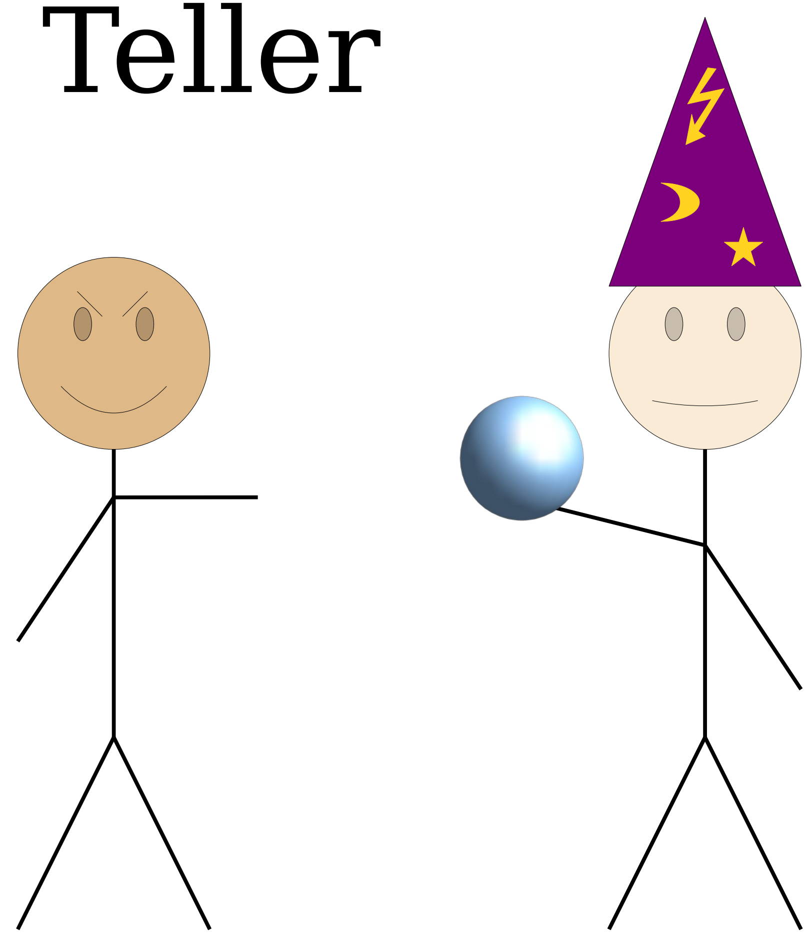
- One day, a trickster arrives. The trickster thinks the fortune teller is lying and can't really see the future.
- The trickster says the following:

"I have a yes/no question about the future. But before I ask my question, let's talk payment.

If you answer yes, then I'll pay you $137.

If you answer no, then I'll pay you $42.

- The fortune teller thinks for a moment, then agrees.



Trickster pays $137 if the fortune teller answers "yes."

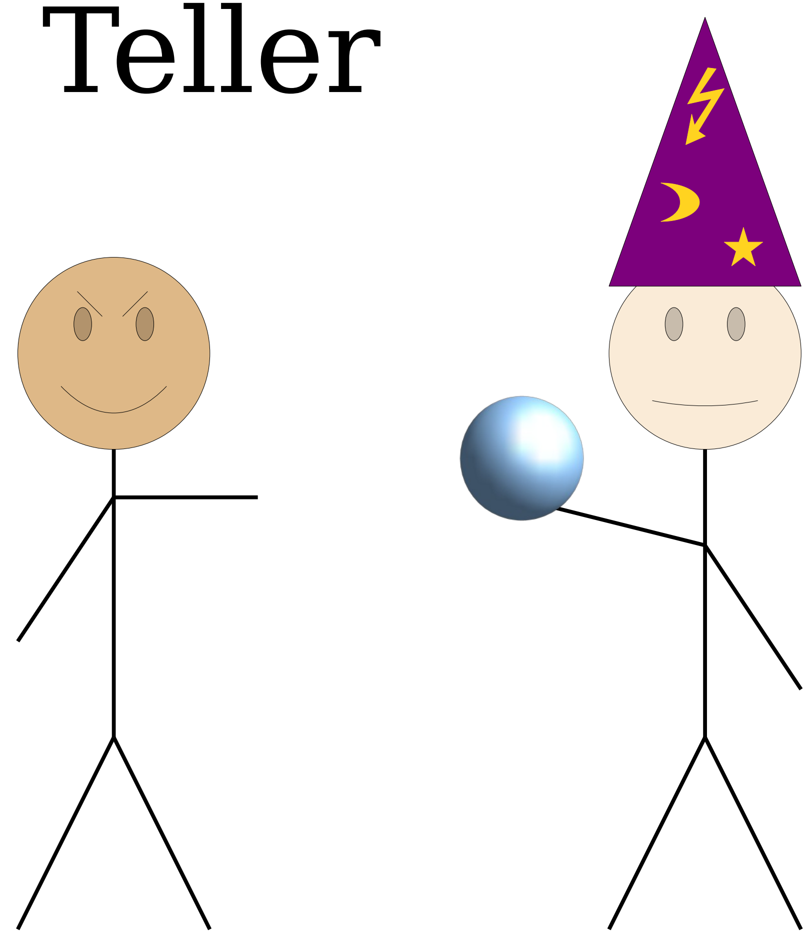
Trickster pays $42 if the fortune teller answers "no."

The Fortune Teller

- The trickster then asks this question:

  "Am I going to pay you $42?"

- The fortune teller is trapped!
- Talk to your neighbor – why?



Trickster pays $137 if the fortune teller answers "yes."

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The Fortune Teller

- The payment scheme the fortune teller agreed to means

  Fortune Teller Says Yes ↔ ***Trickster Pays $137***.

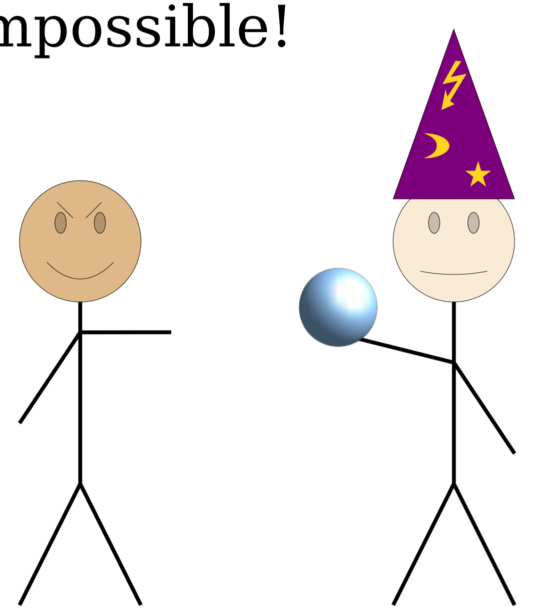
- The trickster's question to the fortune teller means

  Fortune Teller Says Yes ↔ ***Trickster Pays $42***.

- Putting this together, we get

  Trickster Pays $42 ↔ ***Trickster Pays $137***.

- This is impossible!



Trickster pays $137 if the fortune teller answers "yes."

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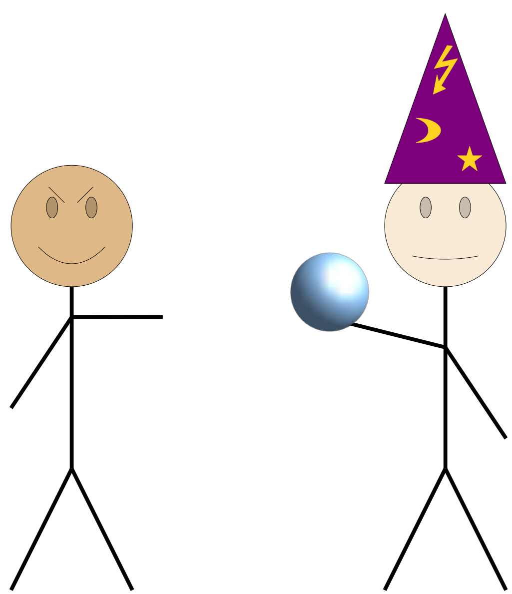
The Fortune Teller

- The fortune teller is a self-defeating object.
- The trickster's strategy is to couple the fortune teller's behavior to what the future holds.
    - The trickster's behavior is chosen in advance to make the fortune teller's answer wrong.
- Therefore, the fortune teller can't answer all questions about all people in the future.



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Part Three: Why Do Programs Loop?

Thoughts on Loops

- In practice, the programs we write sometimes go into infinite loops.
- In Theoryland, Turing machines are allowed to loop. This happens if they don't accept and don't reject.
- ***Question:*** Why are infinite loops possible?
- Or rather: are infinite loops an inherent part of computation, or are they some weird sort of "accident" in how we program computers?

Thoughts on Loops

- ***Theorem:*** The language A_{TM} is recognizable, but undecidable.
  - There's a *recognizer* for A_{TM} (specifically, the universal Turing machine U_{TM}).
  - It is impossible to build a *decider* for this language.
- Stated differently, there's a program we can write (a universal TM) that *has* to loop infinitely on some inputs.
- ***Goal:*** Prove this theorem, and explore its theoretical and philosophical implications.

A_{TM} Revisited

- As a refresher, the language A_{TM} is

  A_{TM} = { ⟨*M*, *w*⟩ | *M* is a TM and *M* accepts *w* }.

- The universal TM U_{TM} has the following behavior when given as input a TM *M* and a string *w*:
  - If *M* accepts *w*, then U_{TM} accepts ⟨*M*, *w*⟩.
  - If *M* rejects *w*, then U_{TM} rejects ⟨*M*, *w*⟩.
  - If *M* loops on *w*, then U_{TM} loops on ⟨*M*, *w*⟩.
- U_{TM} is a recognizer for A_{TM}, but because of that last case it's not a decider for A_{TM}.

A_{TM} Revisited

- As a refresher, the language A_{TM} is

  A_{TM} = { ⟨*M*, *w*⟩ | *M* is a TM and *M* accepts *w* }.

- Given a TM *M* and a string *w*, a decider *D* for A_{TM} would need to have this behavior:
  - If *M* accepts *w*, then *D* **?** ⟨*M*, *w*⟩.
  - If *M* rejects *w*, then *D* **?** ⟨*M*, *w*⟩.
  - If *M* loops on *w*, then *D* **?** ⟨*M*, *w*⟩.
- This is basically the same set of requirements as U_{TM}, except for what happens if *M* loops on *w*.
- Our goal is to prove that there is no way to build a program that meets these requirements.

A_{TM} Revisited

- We can envision a decider for A_{TM} as a function

```
bool willAccept(string fn, string input)
```

  that takes as input the source code of a function (`fn`) and a string representing an input to that function (`input`).
- It then does the following:
  - If `fn(input)` returns true, `willAccept(fn, input)` returns true.
  - If `fn(input)` returns false, `willAccept(fn, input)` returns false.
  - If `fn(input)` loops, then `willAccept(fn, input)` returns false.
- We're going to show it's impossible to write a function that actually does this. But for now, let's just explore what such a decider would do.

```
function = "bool f(string input) {
  if (input == "") return false;
  return input[0] == 'a';
}";

input = "abbababba";

willAccept(function, input) = ?
```

```
function = "bool g(string input) {
  while (true) {
    input += input;
  }
}";

input = "yay! ";

willAccept(function, input) = ?
```

```
function = "bool h(string input) {
  for (char c: input) {
    if (c != input[0]) return true;
  }
  return false;
}";

input = "aaaaaa";

willAccept(function, input) = ?
```

```
function = "bool j(string input) {
  int n = input.length();
  while (n > 1) {
    if (n % 2 == 0) n /= 2;
    else n = 3*n + 1;
  }
  return true;
}";

input = /* 10^137 a's */;

willAccept(function, input) = ?
```

For each of these instances, what does `willAccept(function, input)` return?

Deciding A_{TM}

- Earlier this quarter you explored sums of four squares. Now, let's talk about sums of three cubes.
- Are there integers *x*, *y*, and *z* where…
  - $x^3 + y^3 + z^3 = 10$?
  - $x^3 + y^3 + z^3 = 11$?
  - $x^3 + y^3 + z^3 = 12$?
  - $x^3 + y^3 + z^3 = 13$?

Deciding A_{TM}

- Surprising fact: until 2019, no one knew whether there were integers *x*, *y*, and *z* where

$$x^3 + y^3 + z^3 = 33.$$

- A heavily optimized computer search found this answer:

$$x = 8{,}866{,}128{,}975{,}287{,}528$$
$$y = -8{,}778{,}405{,}442{,}862{,}239$$
$$z = -2{,}736{,}111{,}468{,}807{,}040$$

- As of August 2022, no one knows whether there are integers *x*, *y*, and *z* where

$$x^3 + y^3 + z^3 = 114.$$

Deciding A_{TM}

- Consider the language

$$L = \{ \mathtt{a}^n \mid \exists x \in \mathbb{Z}.\ \exists y \in \mathbb{Z}.\ \exists z \in \mathbb{Z}.\ x^3 + y^3 + z^3 = n \}$$

- Here's code for a recognizer to see whether such a triple exists:

```
bool hasTriple(int n) {
  for (int max = 0; ; max++)
    for (int x = -max; x <= max; x++)
      for (int y = -max; y <= max; y++)
        for (int z = -max; z <= max; z++)
          if (x*x*x + y*y*y + z*z*z == n)
            return true;
}
```

- Imagine calling `willAccept(/* hasTriple code */, 114)`.
  - If such a triple exists, `willAccept` returns true.
  - If no such triple exists, `willAccept` returns false.
- ***Key Intuition:*** However `willAccept` is implemented, it has to be clever enough to resolve open problems in mathematics!

Why is A_{TM} Hard?

- ***Intuition:*** A decider for A_{TM} would be able to…
  - … determine whether the hailstone sequence terminates for any input. (Write a recognizer that runs the hailstone sequence, then feed it into the decider for A_{TM}.)
  - … see if any number is the sum of three cubes. (Write a recognizer that tries all infinitely many triples of integers, then feed it into the decider for A_{TM}.)
  - … and much, much more.
- In other words, this seemingly simple problem of "is this program going to terminate?" accidentally scoops up a bunch of other seemingly harder problems.

Time-Out for Announcements!

Problem Sets

- PS6 was due earlier today. Your diligent CAs are working on grading them! Solution will be released Monday morning
- PS7 has been released and will be due next **Wednesday** at 5:30 PM. Only the coding portion will be mandatory.

Please evaluate this course in Axess.

Your comments really make a difference.

Back to CS103!

Part Four: Self-Referential Software

Self-Referential Programs

- If TMs can take other TMs as input, could they take themselves as input?

YES.

- TMs can take their own code as input, and ask questions about (or even execute!) their own code.
- In fact, any computing system that's equal in power to a Turing machine possesses some mechanism for self-reference.
- Want to see how deep the rabbit hole goes? Take CS154!

Quines

- A ***Quine*** is a special kind of self-referential program that, when run, prints its own source code.
- Believe it or not, it is possible to write such a program!
- *See zip file with lecture slides for code.*

Self-Referential Programs

- ***Claim:*** Going forward, assume that any function has the ability to get access to its own source code.
- This means we can write programs like the one shown here:

```
bool narcissist(string input) {
    string me = /* source code of narcissist */;

    return input == me;
}
```

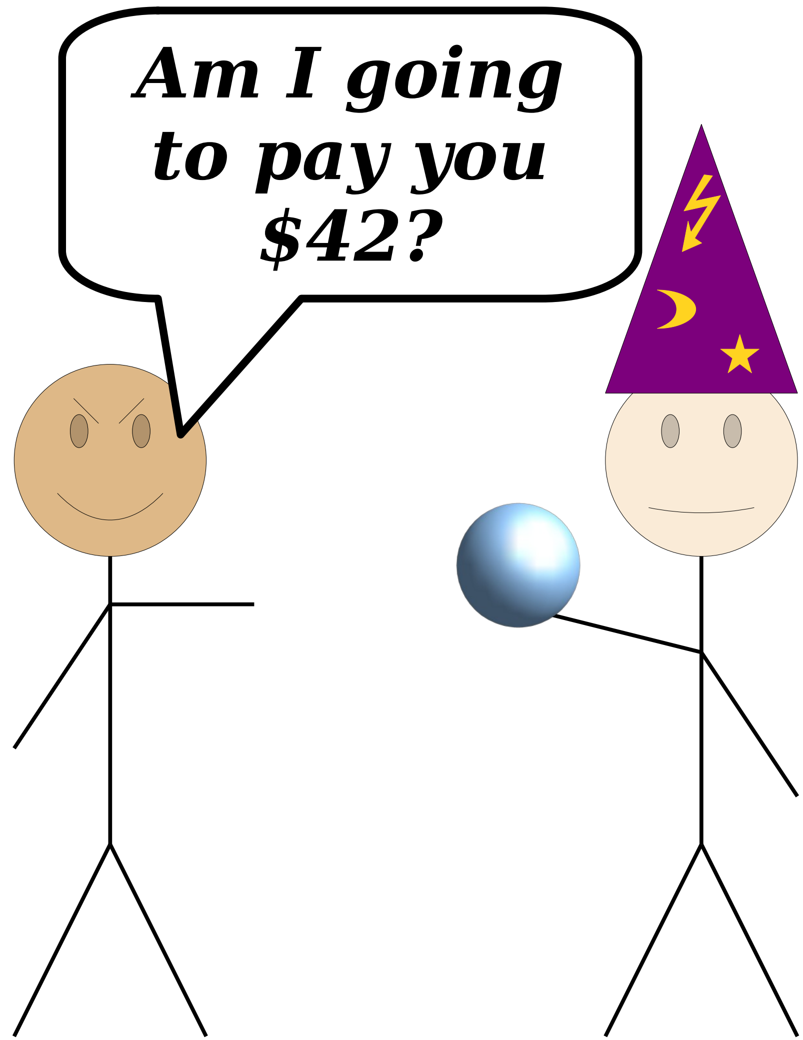
Part Five: Putting It All Together

To Recap

- We're assuming that, somehow, someone wrote a function

```
bool willAccept(string function, string input);
```

  that takes the code of a function and an input to that function, then
    - returns true if `function(input)` returns true, and
    - returns false if `function(input)` doesn't return true.
- ***Goal:*** Show that this decider is "self-defeating;" its power is so great that it undermines itself.
- ***Idea:*** Convert the fortune teller story into a program.



Trickster pays $137 if the fortune teller answers “yes.”

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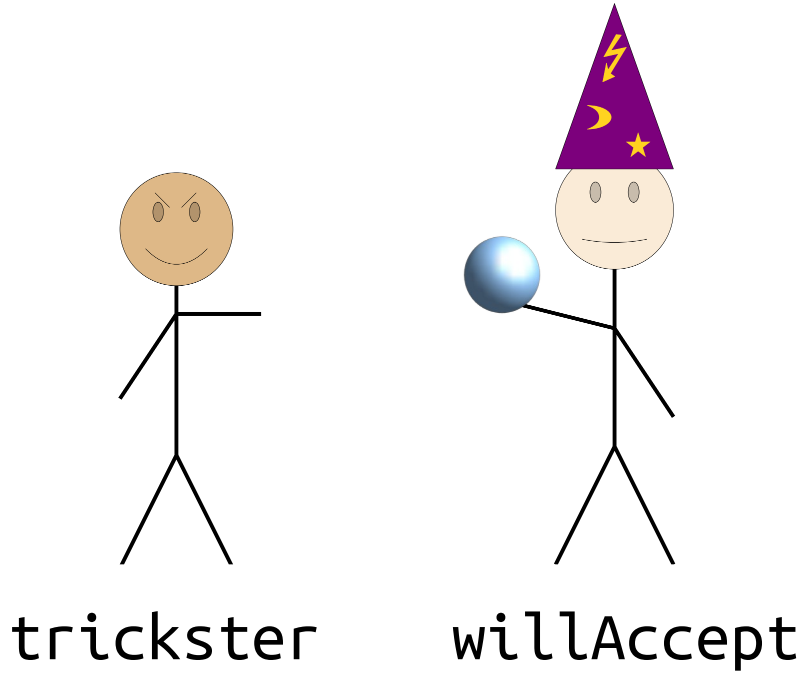
```
bool willAccept(string function, string input) {
   // Returns true if function(input) returns
   // true. Returns false otherwise.
}

bool trickster(string input) {
   string me = /* source code of trickster */;
   return !willAccept(me, input);
}
```



If `willAccept` says `trickster` will return true, then `trickster` returns false.

If `willAccept` says `trickster` will not return true, then `trickster` returns true.

```
bool willAccept(string function, string input) {
    // Returns true if function(input) returns
    // true. Returns false otherwise.
}

bool trickster(string input) {
    string me = /* source code of trickster */;
    return !willAccept(me, input);
}
```

A self-defeating object.

Using that object against itself.

```
bool willAccept(string function, string input) {
    // Returns true if function(input) returns
    // true. Returns false otherwise.
}

bool trickster(string input) {
    string me = /* source code of trickster */;
    return !willAccept(me, input);
}
```

"The largest integer *n*."

"The integer *n* + 1."

Theorem: There is no largest integer.

Proof sketch: Suppose for the sake of contradiction that there is a largest integer. Call that integer *n*.

Consider the integer $n+1$.

Notice that $n < n+1$.

But then *n* isn't the largest integer.

Contradiction! ■-*ish*

Theorem: $A_{TM} \notin \mathbf{R}$.

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Theorem: $A_{TM} \notin \mathbf{R}$.

Proof: By contradiction; assume that $A_{TM} \in \mathbf{R}$. Then there is a decider *D* for A_{TM}. We can represent *D* as a function

```
bool willAccept(string function, string w);
```

that takes in the source code of a function `function` and a string `w`, then returns true if `function(w)` returns true and returns false otherwise.

Theorem: $A_{TM} \notin \mathbf{R}$.

Proof: By contradiction; assume that $A_{TM} \in \mathbf{R}$. Then there is a decider D for A_{TM}. We can represent D as a function

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Given this, consider this function `trickster`:

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Choose a string *w*.

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Choose a string *w*. We consider two cases:

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Case 2: `willAccept(me, input)` returns false.

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Choose a string *w*. We consider two cases:

Case 1: `willAccept(me, input)` returns true. Since `willAccept` decides A_{TM}, this means `trickster(w)` returns true.

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Case 1: `willAccept(me, input)` returns true. Since `willAccept` decides A_{TM}, this means `trickster(w)` returns true. However, given how `trickster` is written, in this case `trickster(w)` returns false.

Case 2: `willAccept(me, input)` returns false.

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Case 2: `willAccept(me, input)` returns false. Since `willAccept` decides A_{TM}, this means `trickster(w)` doesn't return true.

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Case 2: `willAccept(me, input)` returns false. Since `willAccept` decides A_{TM}, this means `trickster(w)` doesn't return true. However, given how `trickster` is written, in this case `trickster(w)` returns true.

Theorem: $A_{TM} \notin \mathbf{R}$.

Proof: By contradiction; assume that $A_{TM} \in \mathbf{R}$. Then there is a decider D for A_{TM}. We can represent D as a function

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In both cases we reach a contradiction, so our assumption must have been wrong.

Theorem: $A_{TM} \notin \mathbf{R}$.

Proof: By contradiction; assume that $A_{TM} \in \mathbf{R}$. Then there is a decider *D* for A_{TM}. We can represent *D* as a function

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In both cases we reach a contradiction, so our assumption must have been wrong. Therefore, $A_{TM} \notin \mathbf{R}$.

Theorem: $A_{TM} \notin \mathbf{R}$.

Proof: By contradiction; assume that $A_{TM} \in \mathbf{R}$. Then there is a decider D for A_{TM}. We can represent D as a function

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```

that takes in the source code of a function `function` and a string `w`, then returns true if `function(w)` returns true and returns false otherwise.

Given this, consider this function `trickster`:

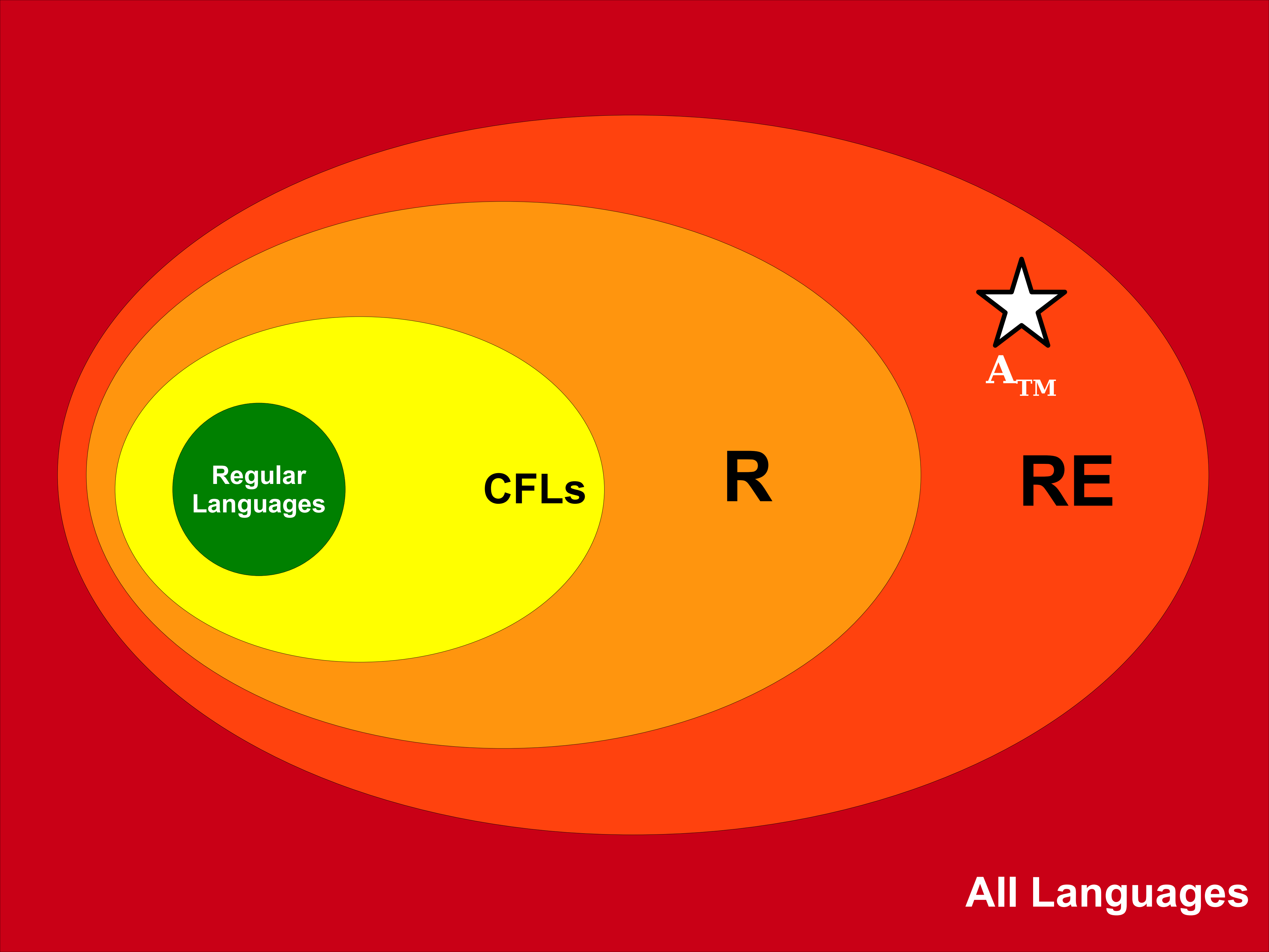
```
bool trickster(string input) {
    string me = /* source code of trickster */;
    return !willAccept(me, input);
}
```

Choose a string w. We consider two cases:

Case 1: `willAccept(me, input)` returns true. Since `willAccept` decides A_{TM}, this means `trickster(w)` returns true. However, given how `trickster` is written, in this case `trickster(w)` returns false.

Case 2: `willAccept(me, input)` returns false. Since `willAccept` decides A_{TM}, this means `trickster(w)` doesn't return true. However, given how `trickster` is written, in this case `trickster(w)` returns true.

In both cases we reach a contradiction, so our assumption must have been wrong. Therefore, $A_{TM} \notin \mathbf{R}$. ■



Regular Languages
CFLs
R
RE
A_{TM}
All Languages

What Does This Mean?

- In one fell swoop, we've proven that
  - A_{TM} is ***undecidable***; there is no general algorithm that can determine whether a TM will accept a string.
  - $\mathbf{R} \neq \mathbf{RE}$, because $A_{TM} \notin \mathbf{R}$ but $A_{TM} \in \mathbf{RE}$.
- What do these three statements really mean? As in, why should you care?

$A_{TM} \notin \mathbf{R}$

- What exactly does it mean for A_{TM} to be undecidable?

  Intuition: The only general way to find out what a program will do is to run it.

- As you'll see, this means that it's provably impossible for computers to be able to answer most questions about what a program will do.

$A_{TM} \notin \mathbf{R}$

- At a more fundamental level, the existence of undecidable problems tells us the following:

  There is a difference between what is true and what we can discover is true.

- Given a TM *M* and a string *w*, one of these two statements is true:

  M* accepts *w ***M* does not accept *w***

  But since A_{TM} is undecidable, there is no algorithm that can always determine which of these statements is true!

R ≠ RE

- Because **R ≠ RE**, there is a difference between decidability and recognizability:

In some sense, it is fundamentally harder to solve a problem than it is to check an answer.

- There are problems where, when the answer is "yes," you can confirm it (run a recognizer), but where if you don't have the answer, you can't come up with it in a mechanical way (build a decider).

Next Time

- ***Why All This Matters***
  - Important, practical, undecidable problems.
- ***Intuiting RE***
  - What exactly is the class **RE** all about?
- ***Verifiers***
  - A totally different perspective on problem solving.
- ***Beyond RE***
  - Finding an impossible problem using very familiar techniques.